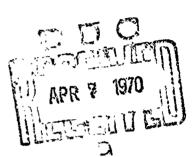
United States Naval Postgraduate School





THESIS

A MODEL FOR RELATING WEAPONS SYSTEM AVAILABILITY AND CONTINUOUS REVIEW INVENTORY POLICIES FOR REPAIR PARTS

by

Stanley Lee Spaulding

October 1969

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by

Stanley Lee Spaulding Lieutenant Colonel, United States Army B.A., University of Texas, 1964

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Author

Approved by:

Thesis Advisor

Chairman, Department of Operations Analysis

Academic Dean

ABSTRACT

A model is developed for relating continuous review inventory policies for repair parts to system availability. The system consists of S identical unit systems, each of which is a series of k-out-of-n structures. Unit system states are zero or one. An optimal cannibalization policy is assumed. Under this assumption the number of unit systems up is always the maximum possible for any given vector of backorders for the N part types in the system. The distribution of backorders under a (Q,r) policy with Poisson demands for each part type is used to derive expressions for system availability as functions of the Q, r vectors. For simplicity it is assumed that order quantities are set by an operating level in terms of days of supply. A numerical technique is presented for finding the vector of reorder points (safety levels) which minimizes the expected cost of on hand inventory subject to one of two alternative availability constraints.

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LIST OF SYMBOLS AND ABBREVIATIONS

$\mathbf{v}_{\mathbf{k}}$	A random variable taking the values zero or one indicating the state of the kth unit system
Z	A random variable indicating the number of operational unit systems
Z*	The value of Z after an optimal cannibalization
S	The total number of unit systems
E(•)	The expected value operator
A _s	Required availability with respect to supply support defined as a required ratio of E(Z) to the total number of unit systems S, given that failed parts are replaced instantaneously if there are spares on hand
P _{min} (k)	Required assurance that at least k unit systems are up, given that failed parts are replaced instantaneously if there are spares on hand
^λ ki	The ith location where a part is installed on the kth unit system
^k	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$
^	$\sum_{k=1}^{S} \Lambda_k = \text{The set of } n \text{\'s} \text{ locations on all unit}$ $k=1 \text{ systems where parts are installed}$
x _{ki}	A random variable taking the values zero or one indicating whether or not there is a working part of the proper type installed in the ith location on the kth unit system.
<u>x</u> _k	A random vector whose components are $X_{k,i}$ $i=1,\ldots,n$ indicating the states of all the locations requiring parts on the kth unit system
<u>x</u>	$(\underline{x}_1,\ldots,\underline{x}_S)$, random vector indicating the states of all locations requiring parts in the whole system
$*_k$	The set of possible values of \underline{x}_k

×	The set of all possible values of \underline{X}
ф	The empty set
FSN j	the jth part type or the jth Federal Stock Number, j = 1,,N
Φ	$\{\hat{rsn}_j\}_{j=1}^N$ = The set of different part types installed on each unit system
N	The number of different types of parts or number of different FSN's installed on each unit system, i.e., N = the size of Φ
G (FSN j)	The subset of each $\Lambda_{\bf k}$ consisting of locations requiring part type j
a _j	The size of G(FSNj), i.e., the number of applications of part type j on each unit system
ť	The minimum number of FSN; which must be working on a unit system if it is to be operational
ť	A random variable indicating the number of working parts of type j in the whole system
$\overline{\mathbf{M}}$	A random vector, (W_1, \dots, W_N) , where W_j is defined above
ť	A random variable indicating the number of backorders for part type j
<u>Y</u> .	A random vector, (Y_1, \dots, Y_N) , where Y_j is \cdots determined above
$1^{K_{\underline{c}}}(\overline{\mathbf{x}})$	The indicator function of the set K at the point \underline{X}
t	Time
τ	Procurement lead time
Q	Order quantity
r	Reorder point
IP	Inventory position
ψ _j (y _j ;r _j)	The probability mass function for backorders for part type j, given that the reorder point is rj

p(u;pt)	The Poisson probability mass function at u_{\star} given that the mean is $\ \rho\tau$
Ρ(u;ρτ)	The complementary cumulative Poisson distribution function at u , given that the mean is $\rho\tau$
ê _j	The failure rate for part type j
j	$L_a Sa_j \hat{\rho}_j$, the pooled failure rate for all parts of type $$ j in the whole system, where L_a is defined below
^L a	Activity level multiplier used to account for varying levels of usage of the system
$\mathbf{r}^{\mathbf{c}}$	Operating level in terms of days or months of supply used to set order quantities

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I. INTRODUCTION

A. BACKGROUND

In the management of logistic support for weapons systems a significant problem in resource allocation is the determination of optimal inventory policies for repair parts. Inadequate stocks of repair parts result in low system availability. Stocks sufficient to insure with high probability that any part will be immediately available when needed my tie-up more resources than are justified.

An indication of the concern of the logistics manager with this problem is the following statement from the revised General Objective No. 1 for the Navy Supply Systems Command: "For technical material, optimum support is that which minimizes downtime of weapons systems due to lack of repair parts and components." [Ref. 1] In order to accomplish this objective, techniques are needed for relating repair parts inventory policies to weapons systems availability.

The mathematical models designed to deal with the impact of repair parts shortages have usually been formulated as a minimization of total cost, where the total cost includes ordering, holding, and backorder costs. In theory the backorder costs are a measure of the impact of shortages on system availability. A difficulty in practice is that backorder costs are hard to estimate. One technique for

finding a value of the shortage cost is to impute it based on a required maximum probability Pout of being out of stock. See Deemer and Hoekstra [Ref. 2, p. 5] for a discussion of this technique. The required value of Pout is determined by management judgement. This judgement presumably includes some intuitive consideration of the effect of repair parts shortages on system availability.

The total cost minimization formulation does not, however, permit any direct correlation of inventory policies and system availability. Inventory policies computed using that formulation are determined for each part independent of all other parts. System availability is, however, a function of $(Q,\underline{r}) = (Q_1,r_1),\ldots,(Q_N,r_N)$. Therefore, if we wish to correlate inventory policies to system availability, we need a multi-item model rather than one which deals with each item one at a time.

There do exist a number of techniques for explicitly correlating system availability and single period inventory policies. One of these is the optimal redundancy approach discussed in Chapter 6 of Barlow and Prochan [Ref. 3]. The problem discussed there is one where at the beginning of the period there is a quantity n; of each part type j on hand, and resupply is not available until the end of the period. The object is to determine the value of each n; so that system availability is maximized subject to some cost constraint or, alternatively, to minimize the cost of achieving a required availability. The cost of the

parts inventory may be interpreted as dollars, weight, volume or some other measure of the amount of a "resource" which is used up by the inventory.

B. GENERAL NATURE OF THE MODEL

The approach taken here is to apply the redundancy optimization ideas to a continuous review inventory situation. In particular the model addresses the problem of optimal parts inventory policies in support of a weapons system at the direct support echelon. The objective function to be minimized is the expected cost of on hand inventory. The constraint is a system availability requirement, which may be in one of two alternative forms. The first is that the expected number of unit systems up must be greater than or equal to a required fraction of the total number of unit systems. The second form of the constraint is that the probability that at least k unit systems are up must be equal to or greater than a required assurance level.

The model is an idealization of the repair parts supply support furnished by the direct support maintenance unit for the population of a particular major item in an operational Army unit. The system consists of S identical unit systems, which might be tanks, aircraft, howitzers, or some other set of equipments of the same make and model.

C. ASSUMPTIONS OF THE MODEL AND THEIR IMPLICATIONS

Before turning to the formal description of the model we shall discuss some of the important assumptions and their implications here.

1. The Inventory System is Single Echelon

To simplify the problem we assume that the supported military unit does not carry a stock of repair parts, but that parts are immediately available from the direct support maintenance unit if there is stock on hand at the time a demand occurs. The inventory system thus is single echelon.

2. Procurement Lead Time is Constant

The procurement lead time (order and ship time) is assumed to be constant for each part type, where part type indicates a particular Federal Stock Number. Lead times may, however, be different for different part types. We further assume that the source of supply is never out of stock. Thus, the model ignores variability of lead times and the possibility of reducing lead times by using a higher requisition priority.

3. Unit System State Values are Zero or One

The supported system consists of S identical unit systems. Each unit system is assumed to exist in one of exactly two states depending upon the states of the installed parts: either it is up, i.e., it is fully capable of operating satisfactorily; or it is who, i.e., totally ineffective. This assumption means that the possibility that the unit system may be partially effective is not

considered. Actual equipment does exist in partially effective states. For example, if the 50-caliber machine gun on a tank is not working, the capability of the tank is reduced in some of its roles, but it is obviously not totally ineffective if all other components of the tank The zero-one assumption is made for mathematiare working. cal simplicity. The two state categorization is not too different from the equipment serviceability code ratings currently used within the Army for determining the materiel readiness of Army units. Possible ESC ratings are green, meaning fully operational and capable of operating in combat for 60 days; amber, meaning fully operational and capable of operating in combat for 30 days; and red, meaning not fully operational. We have simply dropped the amber category.

4. Component Part State Values are Zero or One Similar to the zero-one assumption for each unit system two possible states are assumed for each installed part: either it is working or it has failed. Thus the part is 100 per cent effective or totally ineffective.

5. The Unit System is a Series of k-out-of-n Structures

As mentioned in Section 3 above, the state of a unit system is determined by the states of its component parts. If there are a parts of type j initially installed on each unit system, it is assumed that there is a number b equal to or less than a such that at least b parts of type j must be working if the unit system is

to be up. If the number of working type j parts is less than b_j, the unit system will be down. Further we assume that, if the number of working type j parts is equal to or greater than b_j for all j, then the unit system will be up. These assumptions may be summarized by saying that the unit system is a series of k-out-of-n structures, as defined by Birnbaum, Esary and Saunders [Ref. 4, p. 58].

Our b_j corresponds to these authors' k and our a_j to their n.

6. The Number of Unit Systems Up is the Minimum of the Number Up with Respect to Each Part Type

If Z is the number of unit systems up at some arbitrary time and Z_j is the number which would be up if all non-working parts other than type j were replaced by working parts, then Z is called the state of the system, and Z_j is called the state of the system with respect to part type j. We assume that at any time t the state of the system Z equals the minimum over j of the Z_j .

7. An Optimal Cannibalization Policy is Observed

Perhaps the most significant assumption of the model veloped here is that a policy of optimal cannibalization is observed. By optimal we mean that given any vector $\underline{W} = (W_{.},...,W_{N})$, where W_{j} denotes the number of working parts of type j, after the cannibalization operation the maximum possible number of unit systems will be up. The mathematical structure of the supported system is an adaptation of the structure of the cannibalization model of

Hirsch, Meisner and Boll [Ref. 5]. The reason for using this structure is that, given any vector $\underline{Y} = (Y_1, \dots, Y_N)$, where Y denotes the number of backorders of part type there is a unique number systems up. This would not be the case if cannibalization were not allowed. Consider, for example, a system in which each unit system has two parts of type j installed, and both of these must be working if the unit system is to be up. Suppose there are two backorders for this part type. This means that there are two non-working parts of type j in the system. one of these were on one unit system and the other were on a different unit system, then the number of unit systems down for this part type would be two. If, however, both were on one unit system, then only one unit system would be down for this part type. Under a policy of optimal cannibalization there would always be only one unit system down for this backorder situation. Since the probabilities for the number of systems up used in this model are computed based on the probability distributions of the Y_{ij} , the cannibalization assumption is an essential feature of the model. The degree to which cannibalization is actually used to increase system availability in practice depends upon command policy and the practicality of taking parts from one unit system to make another operational. If the supported unit is spread-out geographically, it may not be feasible to cannibalize. Also the maintenance effort in removing and replacing parts under the cannibalization

operation may overload the capability of the maintenance organization. Nevertheless, some cannibalization may be the best strategy for optimizing system availability when there are constraints on the quantities of repair parts the maintenance unit can carry.

8. The Demand Distribution for Each Part Type is Poisson

The demand distribution for each part type is assumed to be Poisson with mean equal to the product of the failure rate for the part type times the number of parts of that type in the system. The failure rate of each part type is assumed to be the same in all of its applications. Poisson assumption implies that the mean time between failures of a given part type is exponentially distributed. It is true that the exponential distribution may be a poor fit for the mean time between failures where an individual part is installed and replaced with a new part of the same type immediately upon failure, particularly when the part is subject to wearout. Cox [Ref. 6, p. 77] indicates however, that the pooled output of a number of renewal processes tends to have the properties of a Poisson process as the number of renewal processes being pooled gets large. sequently, the Poisson assumption may be a good approximation when the total number of parts of a given type in the system is large. Another implication of the assumption is that the distribution of demands is not affected by the number of unit systems down. This in turn implies that an inoperable unit system continues to generate part failures

even after it has gone down. Errors introduced by this assumption will not be too great if the availability rate is reasonably high or if failure rates are estimated on demand histories over a period when the availability rate was about the same as the required availability rate for the period for which inventory policies are being computed.

9. Demand Rates are Linear with Activity Levels

The parameter in the demand distribution is assumed to be a time rate. As discussed by Soland [Ref. 7, p. 45], the natural parameter of the demand process may be miles driven, rounds fired, or some other measure of the usage or activity level of the system. The device used in the model to account for this fact is an "activity level" multiplier L for the pooled demand rates for each part type. Suppose, for example that rounds fired were the natural parameter for the demand process. Further suppose that a certain number of rounds were fired on the average by each unit system per month during the period when demand data were accumulated for estimating failure rates. Now suppose that during the period for which inventory policies are being computed that the programmed number of rounds per unit system per month is doubled. Then the L would be two, and the failure rate estimates for each part type would be doubled. Use of this device assumes that failure rates are linear with the level of usage. assumption is not in general true, but any more realistic means of handling this problem would complicate the model.

For part types not subject to aging, this assumption is probably fairly good.

10. Parts are Replaced Immediately upon Failure

This model assumes that each part is replaced immediately upon failure with zero replacement time if a spare part is on hand at the time of failure. Also, if backorders exist for a particular part type, the failed part will be installed immediately upon receipt by the maintenance unit. Of course, these assumptions are not true in the "real world." To account for the fact that repair times are finite the following procedure could be used. First estimate the availability which we would expect to achieve if there never were any shortage of repair parts. Call this estimate the availability with respect to maintenance. Call the estimated availability predicted by this model the availability with respect to supply support. Then by taking the product of these two estimates we should have a fair estimate of the overall availability to be expected given the inventory policies and the repair capabilities of the maintenance unit. In any case, the reader should keep in mind that availability values computed with this model need to be corrected for down time due to non-zero repair times.

II. MATHEMATICAL MODEL

A. FORMAL DESCRIPTION OF THE SYSTEM STRUCTURE

The system consists of S identical unit systems. On each unit system k there are n loci where parts are installed. Let the λ_{ki} denote the ith locus on the kth unit system and let the set

$$\Lambda_{k} = \{\lambda_{ki}\}_{i=1}^{n}$$

denote the set of loci on the kth unit system. Let the set

$$\tilde{\Lambda} = \frac{1}{k-1} \Lambda_{k}$$

denote the set of all loci on the S unit systems. Let X_{ki} be a random variable indicating the state of the ith locus on the kth unit system, such that

 x_{ki} = 1 if locus λ_{ki} contains a working part of the type required in that locus,

= 0 if locus λ_{ki} fails to contain a working part of the required type.

Now $\underline{x}_k = (x_{k1}, \dots, x_{kn})$ is a vector of zeros and ones which describes the state of the kth unit system. The possible values of \underline{x}_k correspond to the vertices of the unit cube in Euclidan n-space. Let \underline{x}_k denote the set of all possible values of \underline{x}_k . The size of \underline{x}_k is 2^n . Let the random vector

$$\underline{\mathbf{x}} = (\underline{\mathbf{x}}_1, \dots, \underline{\mathbf{x}}_S)$$

be an nS-component vector indicating the state of the entire system. Let χ denote the set of all possible values of \underline{X} . The size of χ is 2^{nS} . Let the set

$$\Phi = \{FSN_j\}_{j=1}^n$$

represent the set of different part types installed in the n loci of each unit system. Further let $G(FSN_j)$ be the subset of Λ_k which has as elements all those loci requiring part type FSN_j . It is assumed that if locus $\lambda_{ki} \in G(FSN_j)$, then $X_{ki} = 1$ if and only if a working part of type FSN_j is installed in that locus. This assumption implies that different FSN's are not substitutes for one another. Let the size of $G(FSN_j) = a_j$, i.e., the number of applications of FSN_j on each unit system is a_j . Assume that for each j there exists a number b_j such that

- 1. $1 \le b_{i} \le a_{i}$;
- 2. If the number of working parts of type FSN_j installed on unit system k is equal to or greater than b_j, the unit system will not fail due to FSN_j;
- 3. If the number of working parts of type FSN_j installed on unit system k is less than b_j , the unit system will fail due to FSN_j ; and
- 4. b_{i} is independent of b_{i} for all $i \neq j$.

Let the state of the unit system with respect to FSN_j , denoted U_{kj} , be defined as the state in which all loci

not in the set $G(FSN_j)$ have working parts installed. In other words U_{kj} describes the state of the unit system if FSN_j were only part type subject to failure. If W_{kj} is a random variable indicating the number of working parts of type FSN_j installed on unit system k, then

$$U_{kj} = 1 \text{ if } W_{kij} \ge b_{j}$$
,
= 0 otherwise.

Now let the state of the system with respect to FSN_j , so denoted Z_j , be defined as $\sum_{k=1}^{\infty} U_{kj}$. Z_j indicates the number of unit systems which would be operational for a given set of values of W_{kj} , $k=1,\ldots,S$ if the only fallable part type were FSN_j .

For a given value of the system state vector X it may be possible to increase the value of Z by cannibalization. We assume that for any value of X cannibalization will be effected in such a way that the maximum value of Z for that X will be achieved. See Hirsch, Miesner and Boll [Ref. 5, p. 336-342] for a detailed description of the cannibalization operation for the type of structure being discussed here. Let Z* denote the state of the system after an optimal cannibalization, and let Z* denote the state of the system with respect to FSN, under a policy of optimal cannibalization.

Note that Z_j^* may be considered to be a function of W_j the number of working parts of type j, since Z_j^* is

the value of the number of operational unit systems after an optimal cannibalization when the only part failures are at those loci associated with part type j. It is assumed that $Z_j^*(W_j)$ is a monotone non-decreasing function of W_j . Let $K(W_j,z)$ denote the set of values of the number of working parts W_j such that the value of Z_j^* is at least as great as z_j i.e.,

$$K(W_{j},z) = \{W_{j}: 0 \le W_{j} \le Sa_{j} \text{ and } Z_{j}^{*}(W_{j}) \ge z\}.$$

Note that

$$K(W_{j},0) \subseteq K(W_{j},1) \subseteq ... \subseteq K(W_{j},S) \subseteq K(W_{j},S+1), j=1,...,N;$$
 and since max $Z_{j}^{*} = S$,

$$K(W_{\dagger},S+1) = \phi$$
.

The largest value of z such that $K(W_j,z) \neq \varphi$ is S.

Further note that the minimum value of W_j such that Z $Z_j^* = z$ is zb_j since b_j is the minimum number of part type j needed on each unit system if that unit system is to be operational, in symbols

$$zb_{j} = min \{w_{j}: w_{j} \in K(W_{j}, z)\}$$

Let K_{EX} be an arbitrary set where χ is the set of all system state vectors \underline{X} . The indicator function I_K is defined for all K and X as follows

$$I_{K}(\underline{X}) = 1$$
, if $\underline{X} \in K$,
= 0, if $\underline{X} \notin K$.

Now let us use the symbol $\{w_j \geq zb_j\}$ to denote the set of all \underline{X} such that $w_j \geq zb_j$, i.e., $\{\underline{X} \in X \colon w_j(\underline{X}) \geq zb_j\}$. Hirsch, Meisner and Boll show (using somewhat different notation) that if $Z^* = \min_j Z^*_j$, then

$$z^* = \sum_{k=1}^{S} I_{\{w_j^{\geq}zb_j^{}\}} = \sum_{k=1}^{S} \prod_{j=1}^{N} I_{\{w_j^{} \geq zb_j^{}\}}$$

and the authors call this the "representation theorem" [Ref. 5, p. 349]. If S = 1, i.e., we have only one unit system then $Z^* = U$, and U = 1 if and only if each one of the indicators $I_{\{w_j = b_j\}}$ is equal to 1. The structure of the unit system in this model is thus a coherent structure in the definition of that term given by Birnbaum, Esary and Saunders [Ref. 4, p. 61]. Further each unit system may be considered to be a k-out-of-n structure with respect to each part type j, where a k-out-of-n structure is one which has n parts of a single type and is operational if and only if at least k out of the n parts are working. Also, we can consider each unit system to be a series system composed of N k-out-of-n structures, where the k's are the bj's and the n's are the aj's.

B. CORRELATION OF INVENTORY POLICIES AND SYSTEM AVAILABILITY

Consider now the effect of inventory policies on the expected value of the number of operational unit systems E(Z). If there were an infinite number of repair parts of each type available, the value of \underline{X} would always be

 $\underline{1} = (1,1,\ldots,1)$, since it is assumed that each part is replaced immediately upon failure as long as spares are available, and the state of the system Z would always be S. There are in fact various constraints on the ability of the support unit to carry inventories of parts. The result is that from time to time demands will occur for which no replacement part is immediately available. For any state vector \underline{X} , denote the state vector after an optimal cannibalization as \underline{X}^* and the value of Z after cannibalization as \underline{X}^* . Then

$$\underline{X}^{*}(t) = (X_{1,1}^{*}(t), \dots, X_{S,n}^{*}(t)), t \ge 0,$$

is a stochastic process where for each fixed t, X_{ki} is a random variable taking the values zero or one. Now for each subset K of χ , the set of possible values of χ , let $P_t(K)$ denote the probability that $\chi^*(t)$ is an element of the set K, i.e.,

$$P_{+}(k) = Pr\{X*(t) \in K\}, K \in X$$

As shown by Hirsch, Meisner and Boll the probability distribution defined above concentrates all of its mass on the set of maximum points $M_{\underline{X}}$, which are the possible vectors \underline{X}^* , the state vector after an optimal cannibalization. Any vector \underline{X} in which the number of working parts of each type is given by the vector \underline{W} can be transformed into another vector \underline{X}' which has the same value of \underline{W} . Thus, since an optimal cannibalization yields

the maximum value of the system structure function Z for a given state vector \underline{X} , there is a unique value of the state of the system after an optimal cannibalization Z^* for each value of the vector $\underline{W} = (W_1, \dots, W_N)$. The set of all \underline{X} corresponding to a given \underline{W} vector is called an equivalence class, and the set $\underline{M}_{\underline{X}}$ is a subset of this equivalence class where Z as a function of \underline{X} takes its maximum values over the equivalence class. We denote each \underline{X} in $\underline{M}_{\underline{X}}$ as \underline{X}^* . To describe the variation in time of the system state function Z, set

$$Z^*(t) = Z(X^*(t)), t \ge 0,$$

and we note that the probability distribution of Z*(t) is given by

$$Pr{Z*(t) \ge z} = Pr{Z*(x*(t)) \ge z}$$
$$= P_{t}{\underbrace{X: Z*(\underline{X}) \ge z}}.$$

Let W_j (t) represent the number of working spares of type j at time t. Under the assumption that $Z^* = \min_{j \in J} Z^*_j$, we have immediately from the representation theorem that

$$Z^*(t) = \sum_{k=1}^{S} I_{\substack{N \\ j=1}} \{W_j(t) \ge kb_j\} = \sum_{k=1}^{S} \prod_{j=1}^{N} \{W_j(t) \ge kb_j\}.$$

Assume that

1. For each index j = 1,...,N, the loci in which parts of type FSN_j occur are indistinguishable in their

effects on the lifetimes of the parts installed in them, i.e., the failure rate at each instant of a given part type FSN; does not depend on the locus in G(FSN;) in which the part is installed, nor on the particular sequence of loci through which it has passed.

2. Parts operate independently, i.e., the lifetime of a given part is not related to the lifetimes of any other parts. These assumptions make it reasonable to postulate that the joint distribution of $(W_1(t),\ldots,W_N(t))$ does not depend upon the particular cannibalizations involved in the process $\{Z^*(t),\,t\geq 0\}$ and that the N stochastic processes

$$\{W_1(t), t \ge 0\}, \dots, \{W_N(t), t \ge 0\}$$

are mutually independent. Thus the expected number of unit systems up under optimal cannibalization at time t, E(Z*(t)), is given by

$$E(Z^*(t)) = \sum_{k=1}^{S} \prod_{j=1}^{N} E[I_{\{W_{j}(t) \ge kb_{j}\}}]$$

$$= \sum_{k=1}^{S} \prod_{j=1}^{N} Pr\{W_{j}(t) \ge kb_{j}\}.$$

Further, the probability that Z*(t) equals at least z is given by

$$Pr\{Z^*(t) \geq z\} = \prod_{j=1}^{N} Pr\{W_j(t) \geq zb_j\}.$$

Now let Y_j(t) be a random variable denoting the number of backorders for part type FSN_j at time t. Under the assumption that every failed part is replaced immediately by a working part from the inventory of spares, we see that

$$W_{j}(t) = Sa_{j} - Y_{j}(t)$$

and

$$\{W_{j}(t) \ge kb_{j}\} = \{Sa_{j} - Y_{j}(t) \ge kb_{j}\} = \{Y_{j}(t) \le Sa_{j} - kb_{j}\}.$$

Hence

$$Pr\{Z^{k}(t) \geq k\} = \prod_{j=1}^{N} Pr\{Y_{j}(t) \geq Sa_{j} - kb_{j}\}$$

and

$$E(Z^*(t)) = \sum_{k=1}^{S} \sum_{j=1}^{N} Pr\{Y_j(t) \ge Sa_j - kb_j\}.$$

We assume that the inventory system is continuous review with a (Q_j,r_j) policy for each FSN, where Q_j denotes order quantity and r_j denotes reorder point. We also assume the lead time demand is Poisson distributed. Hadley and Whitin [Ref. 8, p. 184] show that under a (Q,r) policy when the lead time demand is Poisson with parameter $\rho\tau$ that the probability that there are y backorders at an arbitrary time t under steady conditions may be derived as follows. The inventory position, IP, is defined as the stock on hand plus stock on order minus backorders. The IP varies between r+1 and r+Q, where r is defined in terms of the IP. Further, Hadley and Whitin show that the

probability that the IP is in any state r+j, $j=1,\ldots,Q$ is 1/Q. Thus if the inventory system is in state r+j at time $t-\tau$, the probability of y backorders at time t equals the probability that y+r+j demands occur during the period $(t-\tau)$ to t for y>0. The probability that y=0 is the probability that the demand during $(t-\tau,t)$ is less than or equal to r+j. Hadley and Whitin do not make this distinction for y=0; consequently, the formula they derive is not valid for y=0. When 0 is substituted for y in the expression they derive the resulting value is the probability that y=0 and on hand inventory = 0. Terms for the probability that y=0 and on hand inventory is greater than zero are left out of their expression. The expression Hadley and Whitin derive is

$$\psi = (1/Q) \sum_{j=1}^{Q} p(y+r+j;\rho\tau) = (1/Q) \sum_{u=y+r+1}^{y+r+Q} p(u;\rho\tau)$$

=
$$(1/Q) [P(y+r+1;\rho\tau) - P(y+r+Q+1;\rho\tau)],$$

where $p(u;\rho\tau)$ is the Poisson probability mass function with parameter $\rho\tau$ at the point u for $u=0,1,2,\ldots$; and $P(u;\rho\tau) = \sum_{v=u}^{\infty} p(v;\rho\tau) \text{ is the Poisson complementary cummulative distribution function.}$ The above expression is valid y greater than zero. For y=0,

$$\psi(0) = (1/Q) \sum_{\Sigma} \sum_{j=1}^{Q} p(u; \rho\tau)$$

=
$$(1/Q)$$
 $\sum_{j=1}^{Q} [1 - P(r+j+1;\rho\tau)]$

= 1 -
$$(1/Q)$$
 $\sum_{j=1}^{Q} P(r+j+1;\rho\tau)$.

We assume that the pooled demand rate $\rho_{\mathbf{j}}$ for all parts of type \mathbf{j} is given by

$$\rho_{j} = L_{a}Sa_{j}\hat{\rho}_{j}$$
,

where L_a is the activity level multiplier discussed in Section I. S is the number of unit systems, a_j is the number of applications of part type j on each unit system, and $\hat{\rho}_j$ is the failure rate for part type j when L_a is one.

Let τ_j be the procurement lead time for part type j. Demand is assumed to be Poisson distributed with parameter $\rho_j \tau_j$. Thus in terms of $\psi_j (y_j)$,

$$Pr\{Z_{j}^{*}(t) \geq k\} = \sum_{\substack{\Sigma \\ Y_{j}=0}}^{Sa_{j}-kb_{j}} \psi_{j}(Y_{j}).$$

Also since $Z^*(t) = \min_{j} Z^*_{j}(t)$ and the y_{j} are assumed to be independent, the probability that $Z^*(t)$ is equal to or greater than k is the probability that each $Z^*_{j}(t)$ is equal to or greater than k, i.e.,

$$\Pr \{\mathbb{Z}^*(t) \geq k\} = \prod_{j=1}^{N} \sum_{y_j=0}^{\operatorname{Sa}_j - kb_j} \psi_j(y_j).$$

Now the expected value of a non-negative integer random variable is the sum from 1 to ∞ of the complementary cumulative probability function; thus

$$E[Z^*(t)] = \sum_{k=1}^{\infty} \{Pr \ Z^*(t) \ge k\} = \sum_{k=1}^{\infty} Pr\{Z^*(t) \ge k\},$$

since the probability that Z*(t) is greater than S is zero. Therefore, the expected value of the number of operational systems in terms of the probability mass functions of the backorders for each part type is:

$$E[Z*(t)] = \sum_{k=1}^{S} \prod_{j=1}^{N} \sum_{y_{i}=0}^{Sa_{j}-kb_{j}} \psi_{j}(y_{j}),$$

where:

$$\psi_{j}(y_{j}) = (1/Q)[P(y_{j}+r_{j}+1) - P(y_{j}+r_{j}+Q_{j}+1)], \text{ for } y_{j}>0,$$

and

$$\psi_{j}(y_{j}) = 1 - (1/Q) \sum_{m=1}^{\bar{Q}} P(r_{j}+m+1), \text{ for } y_{j} = 0.$$

The problem of optimum inventory policies may be formulated as the minimization of the expected cost of on hand inventory subject to one of the following constraints: first, that the expected number of operating systems E(Z) must be equal to or greater than some required fraction of the total number, or, second, that the probability that at least k of the unit systems are operational must be equal to or

greater than a minimum assurance level. The expressions used to relate inventory policies and these constraints are those developed above for E(Z) and $Pr\{Z \ge k\}$. Note that since the expressions refer to "steady state" conditions, and since optimal cannibalization is always performed, the symbol Z will be used henceforth in lieu of $Z^*(t)$ and Z will indicate the "steady state" random variable $Z^*(t)$.

It turns out that for part types with low demand rates the optimal policy under the formulation of the problem indicated above is not to stock these low demand items at the direct support level, but to order them from the source of supply as demands occur. If this is the case, r will be set to -1 and Q to 1. This means that an order is placed as soon as a failure occurs. In this situation the distribution of Y_j reduces to the Poisson distribution with parameter $\rho_j \tau_j$, since the inventory position is always equal to zero, and there is, therefore, only one state for the inventory position. Hence the probability that there are exactly y_j backorders at any arbitrary time is $p(y_j; \rho_j \tau_j)$.

It is interesting to consider what would happen if $\rho_j \tau_j \quad \text{were small for all N part types, and no stock were carried at the direct support level. Assume that each part type were always available from the source of supply with a procurement lead time of <math>\tau_j$, a constant, which might be different for different part types. Now, if no parts are

carried at direct support level, but are ordered only when demands occur, i.e., $r_j = -1$, $Q_j = 1$ for all j, then

$$Pr\{z \ge k\} = \begin{cases} s & saj-kbj \\ ii & s \\ j=1 & y_j=0 \end{cases} p(y_j; \rho_j \tau_j),$$

and.

$$E(Z) = \sum_{k=1}^{S} \prod_{j=1}^{N} \sum_{y_{j}=0}^{Sa_{j}-kb_{j}} p(y_{j}; \rho_{j}\tau_{j}).$$

The values of the above functions can be used to estimate the lower bounds for E(Z) and $Pr\{Z \ge k\}$, when demand rates and procurement lead times are ρ_j, τ_j , $j=1,\ldots,N$.

III. OPTIMIZATION TECHNIQUE AND NUMERICAL ANALYSIS

A. THE COST FUNCTION

The objective function to be minimized is a cost function of the following form:

$$E(C) = \sum_{j=1}^{N} C_{j}D(Q_{j},r_{j}),$$

where

 C_{ij} = The unit cost of part type j,

 $D(Q_j,r_j) = The expected on hand inventory of part type j, given order quantity <math>Q_j$ and reorder point r_j .

If E(C) is in dollars, the above function gives the expected amount of funds tied up in repair parts inventory for the system at the direct support echelon. If E(C) were in pounds the function might represent the expected amount of load carrying capacity of the maintenance used up by the repair parts for the system in question.

The value of $D(Q_{j},r_{j})$ is given by

$$D(Q_{j},r_{j}) = (Q_{j}+1)/2 + r_{j} - \rho_{j}\tau_{j} + B(Q_{j},r_{j}),$$

where

$$B(Q_{j},r_{j}) = \sum_{Y_{j}=0}^{\infty} y_{j}\psi_{j}(Y_{j};Q_{j},r_{j})$$

is the expected number of backorders, given the inventory policy (Q_j, r_j) . Hadley and Whitin [Ref. 8, p. 184] show that

$$B(Q,r) = (1/Q) \sum_{y=0}^{\infty} y[P(y+r+1;\rho\tau) - P(y+r+Q+1;\rho\tau)].$$

B. OPTIMIZATION TECHNIQUE

1. Minimizing Cost Subject to a Required Expected Number of Unit Systems Up

Suppose it is required that the expected number of unit systems up must be equal to or greater than A_sS , where A_s is the required availability with respect to supply support, $0 \le A_s \le 1$, and S is the number of unit systems. Further, suppose it is desired to minimize the cost of expected on hand inventory needed to achieve $E(Z) \ge A_sS$. The problem may then be stated as

$$\begin{array}{ll} \underset{(r_1,\ldots,r_N)}{\text{Minirize}} & \overset{N}{\Sigma} & \overset{C}{C}_{j} D(Q_{j},r_{j}) \\ (q_{1},\ldots,q_{N}) & j=1 \end{array}$$

subject to $E(Z) \ge A_sS$.

Finding the optimal solution to this problem for $(Q_j,r_j)_{ij}$, $j=1,\ldots,N$ is a formidable task due to the complexity of the cost function and the function for the expected value of Z in terms of the Q_j and r_{ji} . To simplify the problem we have chosen to set the values of Q_j by an operating level of supply in terms of days or months of supply. Thus we set

$$Q_{j} = [\rho_{j}L_{o} + .5], \text{ if } [\rho_{j}L_{o} + .5] > 0,$$

= 1, otherwise,

where: [u] = The greatest integer equal to or less than u,

L

The operating level of supply, e.g., one month of supply.

The problem now becomes: given the values of Q_j determined as indicated above, find the vector (r_1,\ldots,r_N) which minimizes the cost function subject to the required level of availability. We shall restrict the r_j to be ≥ -1 . After computing the Q_j values, the next step in the procedure is to compute a table of values of $\psi_j(y_j; r_j = -1)$ for all j. This table is useful in that the values of $\psi_j(y_j; r_j)$ may be found for other values of r_j as follows:

$$\psi_{j}(0; r_{j}) = \sum_{u=0}^{r_{j}+1} \psi_{j}(u; -1), \text{ for } y_{j} = 0,$$

$$\psi_{j}(y_{j}; r_{j}) = \psi_{j}(y_{j}+r_{j}+1; -1), \text{ for } y_{j} > 0.$$

The number of values needed for each ψ_j depend upon the demand rate and the accuracy desired in the computations. After the table of ψ values is computed the next step is to find a set (r_1^*,\ldots,r_N^*) of r_j such that $E(Z_j)$, the expected value of Z with respect to part type j, is equal to or greater than A_SS , and such that $Pr(Z_j=S;r_j)$ is concave for $r_j \geq r_j^*$. Each r_j must be large enough

to insure $E(Z_j) \geq A_s S$ since $E(Z_j)$ gives the expected number of operating systems when every other type except j is always in supply. E(Z) can not be greater than $E(Z_j)$ because it is assumed that at any arbitrary time $Z = \min_{j} Z_j$. Now $E(Z_j)$ is given by

$$E(z_{j}) = \sum_{k=1}^{S} \sum_{y_{j}=0}^{Sa_{j}-kb_{j}} \psi_{j}(y_{j}; r_{j}).$$

Using this expression find the minimum r_j such that $E(Z_j) \geq A_S S$. (See Appendix A for a discussion of the need to require each $\Pr(Z_j \geq k; r_j)$ to be concave for $r_j \geq r_j^*$.) After the set of initial values of the r_j is found, a marginal analysis technique similar to the approach of Karr and Geisler [Ref. 9] and to one form of optimal redundancy algorithm of Barlow and Prochan [Ref. 3, p. 166] is used to find at each iteration the part type which yields the greatest increase in E(Z) for the increase in the cost of the expected on hand inventory for that part type. On the first iteration the ratio

$$\hat{\Delta}_{j} = \frac{E(Z; r_{j}+1) - E(Z; r_{j})}{C_{j}[D_{j}(r_{j}+1) - D_{j}(r_{j})]}$$

is computed for each part type, where

$$E(Z;r_{j}+1) = \sum_{k=1}^{S} Pr(Z \ge k) \frac{Pr(Z_{j} \ge k;r_{j}+1)}{Pr(Z_{j} \ge k;r_{j})}.$$

The maximum over j of Δ_j is determined and the r_j corresponding to that maximum is increased by one. E(Z) is set equal to $E(Z;r_j+1)$ corresponding to this j. If the new value of E(Z) is equal to or greater than A_SS , the procedure stops; if not, a new set of Δ_j is computed and the r_j corresponding to the maximum over j of the new Δ_j is increased by one. The procedure continues in this fashion until the value of E(Z) is equal to or greater than A_SS .

2. Minimizing Cost Subject to a Required Probability that at Least k Unit Systems are Up

Suppose that it is required that $P\{z \ge k\} \ge P_{\min}(k)$, and that it is desired to minimize the cost of the expected on hand inventory needed to achieve this probability. The general procedure for finding the optimum vector (r_1, \ldots, r_N) is the same as the $E(Z) \ge A_SS$ form of the constraint with the following expections:

a. The initial r_j values (r_1^*,\ldots,r_N^*) are the minimum values such that $\Pr(Z_j \geq k) \geq \Pr_{\min}(k)$ and $\Pr(Z_j \geq k_j r_j)$ is concave in r_j for $r_j \geq r_j^*$.

b. At each iteration Δ_j is computed as

$$\Delta_{j} = \frac{\Pr\{Z \geq k; r_{j}+1\} - \Pr\{Z \geq k; r_{j}\}}{C_{j}[D(r_{j}+1) - D(r_{j})]}.$$

1-

As in the first form of the constraint the r_j associated with the maximum Δ_j is increased by one at each iteration until $Pr(Z \ge k) \ge P_{min}(k)$. Since $Pr\{Z \ge k\}$ is just one

term of the sum over k in the expression for E(Z) the computational effort at each iteration is much reduced for this form of the constraint, particularly if the number of unit systems S is large.

C. NUMERICAL ANALYSIS AND DEVELOPMENT OF FORTRAN PROGRAMS

Two FORTRAN programs, one for each form of the constraint have been developed and run for several sample problems on the Naval Postgraduate School's IBM 360 computer. The programs are very similar, and both use identical versions of the following subprograms: PSITAB, which computes a table of values of $\psi_j(y;r_j=-1)$ for $j=1,\ldots,N$; PRZJK, a function subprogram which computes $\Pr\{Z_j \geq k;r_j\}$; EBO, a function which computes the expected number of backorders for part type j give that the reorder point is r_j ; PPT, a function which computes individual terms of the Poisson probability mass function; and PCUMT, which computes complementary cumulative terms of the Poisson distribution.

With the IBM 360 it was found necessary to do the arithmetic in double precision because the single precision round-off errors for repeated multiplications were excessive. The table of $\psi_j(y;r_j=-1)$ values was, however, stored as a single precision array to conserve storage space. When values from this table were needed they were converted to double precision with the standard function DBLE.

Subroutine PSITAB stores the computed values of $\psi_{j}(y;r_{j}=-1)$ in an array called PI(J,K), where

$$PI(J,K) = \Psi_{J}(k-1;r_{J} = -1).$$

The array size is (N,100), where N is the number of different part types; however, for each J only values of PI(J,K) for K less than or equal to KMAX(J) are computed, where

$$KMAX(J) = min \begin{cases} 100, & largest K such that \\ PI(J,K) < 10^{-12} \end{cases}.$$

PI(J,K) for K greater than KMAX(J) are set to zero. For part types with low demand rates the number of values of PI(J,K) computed is much less than 100. The array size used in the programs permits handling demand rates up to about 50 without significant truncation of the distribution of Y_j .

Program EZMIN

Program EZMIN computes the optimum \underline{r} vector subject to a required minimum expected value of the number unit systems up. The main program reads the input data, writes the system parameters and the required availability, and calls subroutines PSITAB, INITAL, and OPTIMR. Subroutine PSITAB computes the PI(J,K) array as discussed above. Subroutines INITAL computes an initial set of r_j values, r_1^*, \ldots, r_N^* , each of which satisfies the following conditions: $\Pr(Z_j = S; r_j)$ a concave function of r_j for $r_j \geq r_j^*$;

$$E(Z_j) = \sum_{k=1}^{S} Pr(Z_j \ge k; r_j^*) \ge A_s S.$$

Subroutine OPTIMR uses the marginal analysis technique previously discussed to find the optimal r vector.

2. Program PKSMIN

Program PKSMIN computes the optimum \underline{r} vector subject to a required assurance that at least k out of S unit systems are operational. This program differs from program EZMIN only in the following respects. PKSMIN Main calls subroutine INTLZ instead of INITAL and subroutine OPTMZ instead of OPTIMR. Subroutine INTLZ finds an initial set of r_j values r_1^*, \ldots, r_N^* such that $Pr(Z_j \ge k; r_j)$ is concave for $r_j \ge r_s^*$ and such that $Pr(Z_j \ge k; r_j^*) \ge P_{\min}$. Subroutine OPTMZ computes the optimum \underline{r} vector subject to the required assurance that Z be equal to or greater than k.

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3. Flow Charts

Flow charts for program EZMIN are presented in Appendix B. Flow charts for program PKSMIN are presented in Appendix C.

IV. EXAMPLE PROBLEMS AND DISCUSSION OF RESULTS

A. DESCRIPTION OF THE EXAMPLE SYSTEM

The example system consists of 50 106-mm recoilles rifles, model M40Al. Several factors made this system a convenient example. First, the total number of repair parts is small enough so that the FORTRAN programs developed for solution of the two alternate formulations of the problem can be used without modification on the IBM 360 G-level computer available at the Naval Postgraduate School. Second, Department of the Army Technical Manual 9-1015-221-35 [Ref. 11] contains data from which failure rates may be estimated. Third, actual demand data for a six month period for 62 weapons used in training heavy weapons infantrymen at Fort Ord, California, were made available to the author by the Fort Ord Post Maintenance Section. Finally, price data were available from the microfilm Army Master Data File Selected Management Data File [Ref. 11].

The parts list in TM 9-1050-221-35 contains 287 different Federal Stock Numbers (FSN). Of these only the 159 FSN which were indicated as being combat essential were used for the sample problems. By limiting consideration to only combat essential parts we insure that some of the assumptions of the model are more nearly satisfied than if all part types were considered. For example, the model assumes that for each part type j there is a positive number b, which is the minimum number of working parts of

type j needed for the unit system to be capable of performing its function. A part would not be combat essential if b; were zero. Thus the assumption that the unit system is a series of k-out-of-n structures and that the possible states are zero or one is more nearly true if the unit system consists only of combat essential parts.

The Federal Stock Number, the number of applications per unit system, and data for estimating failure rates were obtained from columns 2, 4, and 6, respectively of Section 2, Appendix B, Ref. 1.0.

B. SAMPLE PROBLEM 1 - REQUIRED EXPECTED NUMBER OF UNIT SYSTEMS UP CONSTRAINT

The first sample problem was computed for the required expected number of unit systems up form of the constraint.

Table I lists a summary of the input data and results for this problem.

TABLE I

SAMPLE PROBLEM 1 - SUMMARY OF INPUT DATA AND COMPUTED RESULTS

,	NUMBER OF UNIT SYSTEMS	50
	NUMBER OF PART TYPES	159
INPUT	REQUIRED AVAILABILITY	0.95
DATA	OPERATING LEVEL	1.0
_	ACTIVITY LEVEL	1.0
	EXPECTED NUMBER OF OPERATIONAL UNIT SYSTEMS	47.58
COMPUTED RESULTS	EXPECTED COST OF ON HAND INVENTORY	\$418.04
	NUMBER OF ITERATIONS OF MARGINAL ANALYSIS PROCEDURE	295

The above data show that for the system consisting of 50 weapons with 159 part types and a required availability of 0.95 program EZMIN computed an \underline{r} vector with an associated expected cost of on hand inventory of \$418.04. The computed value of E(Z) is 47.58, which indicates that the marginal analysis technique overshot the required expected number of systems up by .08, since the required expected number was $50 \times 0.95 = 47.50$.

Table II lists the input parameters and computed results for each FSN. Parts are listed in Federal Item Identification Number (FIIN) sequence, because they are listed this way on the Army Master Data File [Ref. 11]. The FIIN is the last seven digits of the Federal Stock Number. The data cards were sorted into this sequence to facilitate determination of unit prices. The columns of Table II contain the following: j, the sequence number; FSN, the Federal Stock Number; $A = a_{ij}$, the number of applications of part type j on each unit system; $B = b_{i}$, the minimum number of part type j needed for the unit system to be up; $C = C_{ij}$, the unit cost of part type j; RHOHAT = $\hat{\rho}_{i}$, the failure rate for part type j; RHO = ρ_{i} = $\operatorname{Sa}_{\mathbf{j}}\hat{\rho}_{\mathbf{j}}$, the pooled demand rate for all parts of type \mathbf{j} in the system; TAU = τ_{j} , the procurement lead time for part type j; $Q = Q_{i}$, the order quantity; $R = r_{i}$, the reorder point; $EB = B(Q_j, r_j)$, the expected backorders for part type j; and EOH = $D(Q_{i},r_{i})$, the expected on hand inventory for part type j.

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C. SAMPLE PROBLEM 2 - REQUIRED ASSURANCE AT LEAST K UNIT SYSTEMS UP CONSTRAINT

The second sample problem was computed for the required assurance at least k unit systems are up form of the constraint. Table III lists a summary of the input data and results for this problem.

TABLE III

SAMPLE PROBLEM 2 - SUMMARY OF INPUT DATA AND COMPUTED RESULTS

,	NUMBER OF UNIT SYSTEMS	50
	REQUIRED NUMBER OF UNIT SYSTEMS UP	47
INPUT-	REQUIRED ASSURANCE	0.90
DATA	OPERATING LEVEL	1.0
ĺ	ACTIVITY LEVEL	1.0
	PROBABILITY AT LEAST K UNIT SYSTEMS ARE UP	0.91
COMPUTED	EXPECTED NUMBER OF UNIT SYSTEMS UP	47.46
RESULTS	EXPECTED COST OF ON HAND INVENTORY	\$387.88
	NUMBER OF ITERATIONS OF MARGINAL ANALYSIS PROCEDURE	348

The above data show that for the same system considered in Problem 1 with a required assurance of 0.90 that at least 47 unit systems are up program PKSMIN computed an <u>r</u> vector with an associated expected cost of cn hand inventory of \$387.88. The program overshot the required probability by 0.01 since the computed probability is 0.91 and the requirement was 0.90. After the <u>r</u> vector which met the

constraint was computed the program computed the resulting value of E(Z), the expected number of unit systems up.

It is interesting to note that the expected number of unit systems up very nearly meets the requirement of Problem 1, and the expected cost of on hand inventory is also about the same as for Problem 1. The computer execution time for Problem 2 was about 19 seconds compared to slightly over four minutes for Problem 1.

Table IV lists the input paramaters and computed results for each FSN. The column headings are the same as for Problem 1.

- D. EXTECTED COST OF ON HAND INVENTORY AS A FUNCTION OF REQUIRED AVAILABILITY
 - 1. Expected Cost Versus Required Expected Number of Unit Systems Up

Program EZMIN was run for ten values of $A_s = 0.90$, ...,0.99 in order to develop a functional relationship between the first form of the availability requirement and the expected cost of on hand inventory. Table V summaries the results. In addition to the expected cost of on hand inventory the table lists the number of iterations of the marginal analysis procedure required to find a solution and the number of reorder points which were less than the mean lead time demand for each level of A_s . A reorder point less than the mean lead time demand is equivalent to a negative safety level. For many of these the computed order quantity is one and the computed reorder point is minus one.

TABLE IV

2.3066 0.0823 2.6182 8.9602 8.9002 4.9685 5.9064 2.3066 0.0 0.2466 1.3408 0.4966 4.2169 4.2169 1.3408 06.9 8.0 0 SAMPLE PROBLEM 2 - REQUIRED ASSURANCE AT LEAST K UNIT SYSTEMS UP CONSTRAINT INPUT PARAMETERS AND COMPUTED VALUES OF ORDER QUANTITIES, REORDER POINTS, EXPECTED BACKORDERS AND EXPECTED ON HAND QUANTITIES FOR ALL PART TYPES 0.0066 0.0066 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 **EB0** \simeq HEHMAT TO HEHMEHMEHMEN TAU 0.700 0.400 0.700 0.450 0.450 0.450 1.400 3.600 0.800 3.600 0.700 0.400 0.400 0.200 1.400 0.700 6.700 5.600 800 0.01400 0.01500 0.00800 0.00700 0.07200 0.07200 0.01600 0.01400 00800 01400 01400 01400 05600 01400 00200 P.HOHAT 00400 01600 000 10.05 10 ∞ 9135 98135 98138 98138 98138 98138 98147 98147 98188 219 219 219 225 269 53300 53000 53

0.8384 5.2328 5.1580 0.0 0.0 0.3384 0.3612 0.3612 0.3639 0.3639 0.06703 0.06703 0.0703 EBO - CONTINUED 01400 002800 00800 00800 002800 RHOHAT 2 TABLE ပ Q H H H H H H H H H H Q H H 9 9 H H Q Q H H Q H H Z Z H Q F ジェーエーエー チェーーシュー 0 らりょうとりょうこう HSTiiioi NS:

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This (Q,r) implies that the part is not stocked by the maintenance unit but is ordered from the source of supply as demands occur.

The data show that expected costs of on hand inventory rise very steeply as the required level of $\mathbf{A}_{\mathbf{S}}$ approaches unity.

TABLE V ${\tt EXPECTED} {\tt COST} {\tt OF} {\tt ON} {\tt HAND} {\tt INVENTORY} {\tt VERSUS} {\tt REQUIRED} {\tt A}_{{\tt S}}$

As	E (C)	No. of Iterations	No. of Negative Safety Levels
0.90	\$16.80	31	156
0.91	\$22.02	64	151
0.92	\$38,14	67	143
0.93	\$79.57	140	123
0.94	\$168.64	141	104
0.95	\$418.04	295	73
0.96	\$645.76	238	41
0.97	\$1225.36	359	19
0.98	\$1619.00	322	7 `
0.99	\$2406.83	386	2

Figure 1 plots the expected cost of on hand inventory $E\left(C\right)$ as a function of A_{s} .

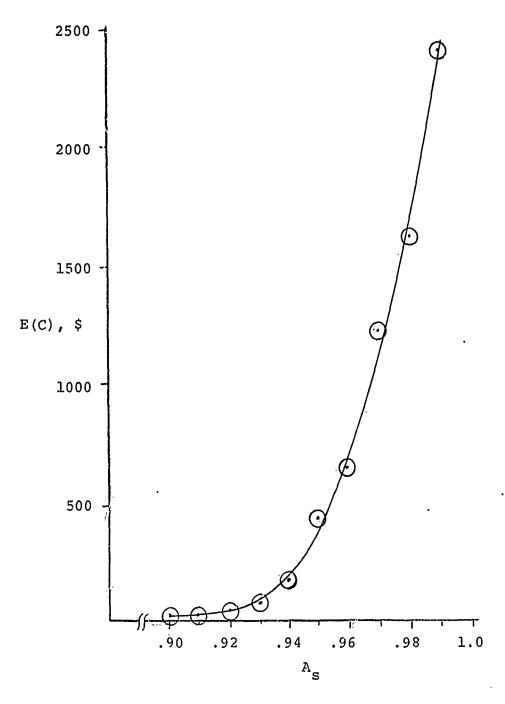


Figure 1. Expected Cost of on Hand Inventory E(C) as a Function of Required Expected Fraction of Unit Systems Up

2. Expected Cost of On Hand Inventory Versus Required Number of Unit Systems Up with an Assurance of 90 Per Cent

Program PKSMIN was run for six values of k=45, ...,50 at a required assurance of 90 per cent in order to develop a functional relationship between the required value of k and the expected cost of on hand inventory at this assurance level. Table VI summaries the results.

TABLE VI

EXPECTED COST OF ON HAND INVENTORY VERSUS REQUIRED NUMBER UNIT SYSTEMS UP WITH 90 PER CENT ASSURANCE

k	E (C)	No. of Iterations	No. of Negative Safety Levels
45	\$29.07	103	150
46	\$84.97	102	122
47	\$387.88	177	85
48	\$1117.47	218	49
49	\$2235.32	278	11
50	\$3519.92	303	0

Figure 2 plots the expected cost of on hand inventory E(C) as a function of the required number k of unit systems up with 90 per cent assurance.

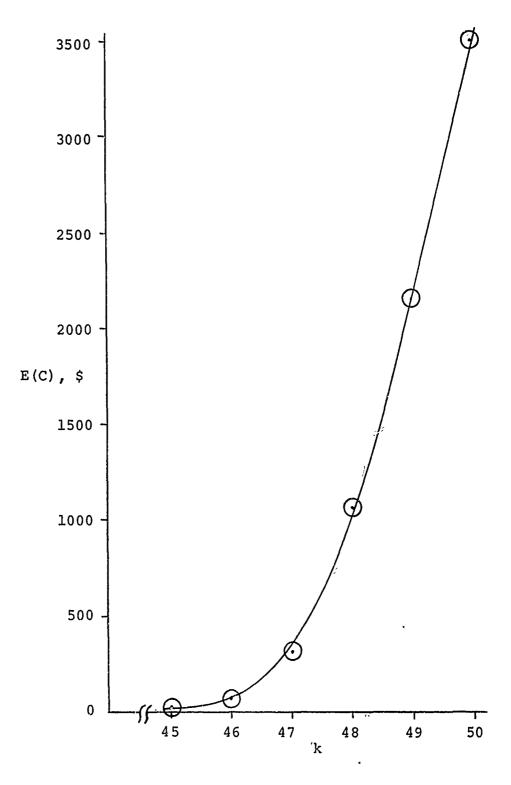


Figure 2. Expected Cost of on Hand Inventory E(C) as a Function of the Required Number k of Unit Systems Up with 90 Per Cent Assurance

V. OBSERVATIONS AND CONCLUSIONS

A. OBSERVATIONS

1. Marginal Analysis may Produce only a Near-Optimal Solution

The marginal analysis procedure presented in Section III sometimes overshoots the optimum solution. That is, the r vector found using the procedure sometimes results in a higher than required availability and a higher expected cost of on hand inventory than needed to satisfy the constraint. This behavior of marginal analysis was observed during the hand calculation of the solution to a small problem with four unit systems and four part types where the constraint was $E(Z) \geq 3.6$. Marginal analysis yielded an E(Z) of 3.709 and an E(C) of \$53.87. was possible, however, on the last iteration to achieve an E(Z) of 3.601 at a cost of \$48.63 by increasing a different r, than the one associated with the maximum A dynamic programming procedure would overcome this shortcoming, but dynamic programming rapidly becomes computationally cumbersome as the number of decision variables In the sample problems discussed in Section IV there were 159 decision variables, and in many possible applications of the model the number would be in the Thus the simple marginal analysis procedure may thousands. be more computationally feasible than dynamic programming for many applications of interest. Further, the relative

differences between the optimal solution and the marginal analysis solution tends to decrease as the size of the problem increases.

2. Program PKSMIN is Much Faster than Program EZMIN

Sample problems 1 and 2 in Section IV resulted in about the same values of E(Z) and of E(C). For problem 1 E(Z) was 47.58 and E(C) was \$418.04. For problem 2 E(Z) was 47.46 and E(C) was \$387.88. On the IBM 360 program EZMIN required over four minutes to compute the results for problem 1, while program PKSMIN solved problem 2 in nineteen seconds. The reason for this difference is that E(Z) a more complex function than $Pr(Z \ge k)$. For a given \underline{r} vector E(Z) is the sum over k from one to S of $Pr(Z \ge k)$. A comparison of the JMAX segments of subroutines OPTIMR and OPTMZ shows that OPTIMR requires more computations at each iteration (See Flow charts B4a and C3a.)

3. Functional Relationships of Expected Cost Versus Required Availability are Readily Obtained

The functional relationships of expected cost of on hand inventory versus availability discussed in Section IV were obtained by simply adding a DO loop to the two main programs. Thus for E(C) versus A_S the value of A_S was initially set to 0.89, and at each iteration of the DO loop A_S was increased by 0.01 for ten iterations. Similarly for E(C) versus k at the 90 per cent assurance level k was initially set at 44 and incremented by one at each iterations for six iterations.

Using the functional relationships of E(C) versus availability we can estimate the maximum availability which can be achieved if there is a constraint on E(C). Suppose, for example, we desired to maximize A_S subject to E(C) \leq \$1000.00. From Table V in Section IV we see that the maximum A_S would be between .96 and .97. To determine maximum A_S more precisely and to determine the associated r vector we could run program EZMIN for a series of values of A_S , say .961, .962, ..., .969. The resulting E(C) for one of these A_S values should be very close to \$1000.00.

Computing time in developing the E(C) versus availability functions can be conserved if at each iteration after the first the starting values of the r_j are the final values computed on the previous iteration. This can be accomplished if the initializing subroutines INITAL and INTLZ are skipped on the second and subsequent iterations. This technique reduces the computing time to about half that needed if the initializing subroutines are called at each iteration.

B. CONCLUSIONS

1. Advantages of the Model

The model provides the first technique of which the author is aware for explicitly taking into account system availability in the determination of continuous review inventory policies for repair parts. The model allows for variable safety levels and yields a set policies

where the ratio of marginal availability to marginal cost of inventory investment is approximately the same for all part types. The result is that the model tends to favor low-cost, high-demand items with relatively high safety levels. High-cost, low-demand items are less well protected. All items are, however, well enough protected to insure, under the assumptions of the model, that the availability requirement will be met.

Compared to models which might more realistically represent the "real world" the one presented here possesses the advantage of computational feasibility. Programs EXMIN and PKSMIN produce solutions in a reasonable amount of time. Refinements in the programming could probably improve run times.

2. Limitations of the Model

The limitations stem primarily from the assumptions which were made for the sake of mathematical simplicity.

For example, the constant lead time assumption is not usually true. Lead times can have considerable variability. Lead time variability will tend to decrease the actual availabilities achieved compared to those predicted by the model. On the other hand the model does not consider the fact that requisitions for parts which are causing a unit system to be down have a higher priority and thus a shorter lead time than normal replenishment requisitions. This fact will tend to make achieved availabilities higher than those predicted.

The optimal cannibalization assumption is seldom completely true. The significance of this assumption decreases, however, as the required availability increases, since at higher availability levels we expect fewer back-orders and therefore fewer occasions for cannibalization.

The significance of the assumption of zero or one states for unit systems will depend upon the stringency of the serviceability criteria used to determine whether a unit system is up or down. The more stringent the criteria, the more parts will have b; equal to a; Parts with b; equal to a; will generally be better protected than parts with b; less than a; The most stringent criteria possible would be that every part on the unit system must be working if it is to be counted as being up. Thus we can see that the inventory policies computed by this model and the resulting expected costs of on hand inventory are highly dependent upon the serviceability criteria for the unit system.

The assumption that parts are replaced immediately with zero lead time makes it necessary to correct the availabilities computed by the model for down time due to finite time to repair.

3. Uses of the Model

The model was developed for use in determining direct support inventory policies for repaire parts for a given population of a particular major item.

Another use which suggests itself is determining optimal maintenance floats. Suppose it were required that at least k unit systems be up with probability .90. Assume that the capability of the maintenance unit to make repairs, given that needed parts are available, is fixed. Then the probability that k unit systems are up is determined by the total number S of unit systems and the inventory policies for repair parts. The quantity S - k is called the maintenance float. An interesting problem is determining the optimal level of the maintenance The model presented in this thesis could be used to estimate the optimal float level as follows. Compute the expected cost of on hand inventory needed to assure with probability .90 that k unit systems are up with zero, one, two, ... unit systems in the maintenance float. Then choose that level of maintenance float which minimizes the sum of the cost of the float plus the on hand inventory.

4. Extensions

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A number of possible extensions of the model appear to be worthy of investigation. For example, relaxation of the assumption that all unit systems are identical would be useful. A next higher level of complexity of the system would be one in which there are two types of unit systems. For example, one type of unit system might be a particular type of weapon. The other unit system might be a piece of fire control equipment which controls the fires of several weapons. If a fire control unit is down all the

weapons it controls are also down. Application of the approach of this model to optimal inventory policies for such a system seems to offer interesting possibilities.

In any case it is hoped that the model presented here will prove useful in the development of system availability-oriented operating policies for military repair parts inventories.

APPENDIX A

CONVENXITY CONSIDERATIONS

In order to insure the optimality of the \underline{r} vector found using the procedure outlined in Section III it is necessary that the cost function be convex and that the region in r-space defined by the constraint be convex.

1. The Cost Function

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The objective function in the model presented in this paper is the expected cost of on hand inventory, which is given by

$$D(C) = \sum_{j=1}^{N} C_{j}D(Q_{j},r_{j}).$$

Since the Q. are fixed by the operating level of supply policy as described in Section III, we are concerned only with convexity with respect to the r_j , $j=1,\ldots,N$. Further, since E(C) is a sum of terms each of which is a function only of r_j , E(C) is convex if each $D_j(Q_j,r_j)$ is convex. Since Q_j is fixed let us drop it as an argument; also let us fix j and drop the subscript. Hadley and Whitin [Ref. 8, p. 184] show that

$$D(r) = (Q + 1)/2 + r + B(Q,r),$$

where B(Q,r) is the expected backorders, given a (Q,r) policy. D(r) is linear in r except for the B(Q,r) term.

Thus D(r) is convex if B(Q,r) is convex in r. Ayain, since Q is fixed, let us drop it as an argument. Hadley and Whitin [Ref. 8, p. 184] show that

$$B(r) = (1/Q) \sum_{y=0}^{\infty} y[P(y+r+1;\mu) - P(y+r+Q+1;\mu)].$$

Let

$$\Delta_{r}B(r) = B(r+1) - B(r)_{r}$$

and

$$\Delta_r^2 B(r) = \Delta_r B(r+1) - \Delta_r B(r)$$
.

D(r) is convex if $\Delta_r^2 B(r) \leq 0$ for all r. Since the parameter μ is fixed for fixed j, let us further simplify the notation by letting

$$P(x) = P(x; \mu)$$
,

$$p(x) = p(x;\mu)$$
.

Further, let

$$B*(r) = QB(r)$$
.

Then D(r) is convex for $r \ge r^*$ if $\Delta_r^2 B^*(x) \ge 0$ for all $r \ge r^*$. Now

$$\Delta_{r}B^{*}(r) = \sum_{y=0}^{\infty} y \begin{bmatrix} P(y+r+2) - P(y+r+Q+2) \\ -P(y+r+1) + P(y+r+Q+1) \end{bmatrix}$$
$$= \sum_{y=1}^{\infty} y [p(y+r+Q+1) - p(y+r+1)].$$

Note that we have dropped the term for y = 0 since it contributes nothing to the sum.

Now

$$\Delta_{r}^{2}B^{*}(r) = \sum_{y=1}^{\infty} y \begin{bmatrix} p(y+r+Q+2) - p(y+r+2) \\ -p(y+r+Q+1) + P(y+r+1) \end{bmatrix}$$
$$= \sum_{y=1}^{\infty} y [\Delta p(y+r+Q+1) - \Delta p(y+r+1)]$$

where

1

$$\Delta p(x) = p(x+1) - p(x)$$

$$= \frac{\mu(x+1) e^{-\mu}}{(x+1)!} - \frac{\mu^{x} e^{-\mu}}{x!}$$

$$= \left[\frac{u - x - 1}{x + 1}\right] p(x).$$

Thus

$$\Delta_{r}^{2}B^{*}(r) = \sum_{y=1}^{\infty} y \left[\frac{y-y-r-Q-2}{y+r+Q+2} p \left(y+r+Q+1\right) - \frac{y-y-r-2}{y+r+2} p \left(y+r+1\right) \right]$$

$$= \sum_{y=1}^{\infty} y \left[\frac{y+r+2-\mu}{y+r+2} p \left(y+r+1\right) - \frac{y+r+Q+2-\mu}{y+r+Q+2} p \left(y+r+Q+2\right) \right]$$

$$= \frac{\left(r+3-\mu\right)}{r+3} p \left(r+2\right) - \frac{\left(r+Q+3-\mu\right)}{r+Q+3} p \left(r+Q+2\right)$$

$$+ \frac{2\left(r+4-\mu\right)}{r+4} p \left(r+3\right) - \frac{2\left(r+Q+4-\mu\right)}{r+Q+4} p \left(r+Q+3\right)$$

$$\vdots$$

$$+ \left(Q+1\right) \frac{\left(r+Q+3\right)}{r+Q+3} p \left(r+Q+2\right) - \dots$$

$$= \sum_{y=1}^{Q-1} y(\frac{r+y+2-\mu}{r+y+2}) p(y+r+1) + Q \sum_{y=Q}^{\infty} (\frac{r+y+2-\mu}{r+y+2}) p(r+y+1).$$

Now every term in this series is non-negative if $r \ge \mu-3$. Thus D(r) is convex in r for $r = \mu-3$. Furthermore, numerical calculations show that for values of μ up to 40 that if $Q \ge \mu-.5$ then D(r) is convex for $r \ge -1$. This numerical result is probably due to the fact that the mode of the Poisson probability mass function is approximately μ -1, which means that the most negative values of (r+y+2-u)/(r+y+2) in the series are multiplied by small values of $p(x;\mu)$. When $r \ge -1$, all terms in the series for $y \ge \mu - 1$ will be positive. Thus it is likely that the sum of the positive terms in the series is greater than the absolute value of the sum of the terms for which $y + 1 - \mu$ is negative.

Thus, as far as the cost function is concerned, we can be sure that the function is convex in \underline{r} for all $\underline{r} \geq \underline{\mu} - \underline{3}$, where $\underline{\mu} = \mu_1, \dots, \mu_N$, and we have reason to believe the function is convex for $\underline{r} \geq -\underline{1}$ when $Q_j \geq \mu_j - .5$, $j = 1, \dots, N$.

2. The Constraint Function

The two alternate forms of the constraint functions are, first, a required expected number of operational unit systems, expressed by

and second a required probability that at least k unit systems are operational, expressed by

$$\prod_{j=1}^{N} \sum_{y_{j}=0}^{Sa_{j}-kb_{j}} \Psi_{j}(y;r_{j}) = P_{\min}(k) .$$

If the constraint region is to be convex, the above functions must be concave functions of $\underline{r}=(r_1,\ldots,r_N)$. Note that the left hand side of the first constraint is the sum over k of terms of the same form as the left hand side of the second constraint. Thus if we can show the conditions for concavity of the left hand side of the second constraint for a general k, we can easily show the conditions for concavity of the first constraint.

The difficulty in showing the concavity of the second form of constraint is that it is not a separable function of the r_j . It can be transformed into a separable function by taking the logarithm, in which case it becomes a sum of terms, each of which is a function of only one r_j , i.e.,

$$ln[Pr{Z \ge k}] = \sum_{j=1}^{N} ln[Pr{Z_j \ge k}],$$

where

$$Pr\{Z_{j} \geq k\} = \sum_{\substack{\sum j=0 \\ y_{j}=0}}^{Sa_{j}-kb_{j}} \Psi_{j}(y_{j};r_{j}).$$

Let

$$f(\underline{r}) = Pr\{z \ge k; \underline{r}\}, = \iint_{j=1}^{N} f_{j}(r_{j}),$$
 $f_{j}(r_{j}) = Pr\{z_{j} \ge k; r_{j}\},$
 $\phi(\underline{r}) = \ln f(\underline{r}),$
 $\phi_{j}(r_{j}) = \ln f_{j}(r_{j}).$

Now it can easily be shown that if $f_j(r_j)$ is concave, $\phi_j(r_j)$ is also concave, and thus $\phi(\underline{r})$ is concave. Unfortunately, the concavity of $\phi(\underline{r})$ does not imply the concavity of $f(\underline{r})$. Now consider

$$f_{j}(r_{j}) = \sum_{\substack{\Sigma \\ y_{j}=0}}^{Sa_{j}-kb_{j}} \psi_{j}(y_{j};r_{j}).$$

Let

$$M = Sa_{j} - kb_{j},$$

then

$$f_{j}(r_{j}) = \Psi_{j}(0;r_{j}) + ... + \Psi_{j}(M;r_{j}).$$

In section III we showed that

$$\Psi_{j}(0;r_{j}) = \sum_{\substack{\Sigma \\ Y_{j}=0}}^{r_{j}+1} \Psi_{j}(Y_{j};-1),$$

$$\Psi_{j}(y_{j};r_{j}) = \Psi_{j}(y_{j} + r_{j} + 1; -1), y_{j} > 0.$$

Thus we have

$$f_{j}(r_{j}) = \sum_{u=0}^{r_{j}+1} \Psi_{j}(u;-1) + \sum_{y=1}^{M} \Psi_{j}(y_{j}+r_{j}+1;-1)$$

$$= \sum_{u=0}^{M+r_{j}+1} \Psi_{j}(u;-1).$$

Now let

$$\Delta_{\hat{r}_{j}}^{f_{j}(r_{j})} = f_{j}(r_{j}+1) - f_{j}(r_{j}),$$

then

$$\Delta_{r_{j}}f(r_{j}) = \sum_{u=0}^{M+r_{j}+2} \Psi_{j}(u;-1) - \sum_{u=0}^{M+r_{j}+1} \Psi_{j}(u;-1)$$
$$= \Psi_{j}(M+r_{j}+2;-1).$$

Let

$$\Delta_{r_{j}}^{2}f_{j} = \Delta_{r_{j}}f(r_{j}+1) - \Delta_{r_{j}}f(r_{j}).$$

then

$$\Delta_{r_{j}}^{2}f_{j} = \Psi_{j}(M+r+3;-1) - \Psi_{j}(M+r+2;-1)$$

$$= (1/Q) \begin{bmatrix} P(M+r_{j}+3) - P(M+Q_{j}+r_{j}+3) \\ -P(M+r+2) - P(M+Q+r+2) \end{bmatrix}$$

$$= (1/Q) [p(M+Q+r+2) - p(M+r+2)]$$

Now f_j is concave for all $r_j \ge r_j^*$ if the expression in brackets above is ≤ 0 for all $r_j \ge r_j^*$. Thus we seek r_j^* such that for all $r_j \ge r_j^*$

$$p(M+r_j+2) \ge p(M+Q_j+r_j+2)$$

or

$$\frac{\mu_{j}^{(M+r_{j}+2)}e^{-\mu}}{(M+r_{j}+2)!} \geq \frac{\mu_{j}^{(M+Q_{j}+r_{j}+2)}e^{-\mu}}{(M+Q_{j}+r_{j}+2)!}$$

or

$$\mu^{Q_{j}} \leq \frac{(M+r_{j}+Q_{j}+2)!}{(M+r_{j}+2)!}.$$

The right hand side of the above inequality is an increasing function of r_j , Q_j , and M. It is usually true that the operating level policy used within the Army is at least as many days of supply as the procurement lead time. This means that $Q_j \geq u_j$. Assume this is true. Now if k = S and $a_j = b_j$, then M = 0. If $f_j(r_j; k=S)$ is concave, $f_j(r_j; k<S)$ will also be concave, since if the inequality holds for M = 0, it will hold for M>0. Let M = 0. Then the inequality becomes

$$\mu^{Q_{j}} \leq \frac{(r_{j}+Q_{j}+2)!}{(r_{j}+2)!}$$

For example, if $r_j = -1$ and $Q_j = 1$, then μ_j must be equal to or less than 2. In other words, for $u_j \le 2$, $f_j(r_j)$ is concave in r_j for all $r_j \ge -1$, $Q_j \ge 1$, and $k \le S$.

To insure the concavity of $f_j(r_j)$ for $\mu > 2$, in the FORTRAN program for the $E(Z) = A_S S$ form of the constraint,

the subroutine INITAL finds an initial set of r_j such that the concavity conditions for k=S are satisfied for r_j equal to or greater than the initial r_j . Similarly, in the FORTRAN program for the $Pr(Z=k) = P_{min}$ form of the constraint, subroutine INTLZ finds an initial set of r_j such that each $Pr(Z_j = k; r_j)$ is a concave function of r_j for r_j equal to or greater than the initial value.

The solution procedure employed in the computer programs thus insures that each $f_j(r_j)$ is concave in r_j in the region that the marginal analysis procedure searches for an optimum. The question of the concavity of $f(\underline{r}) = \prod_{j=1}^{n} f_j(r_j)$ remains open. To date the numerical results have not inidecated that there is a problem, i.e., the numerical results are quite reasonable. No proof, however, of the concavity of $f(\underline{r})$ has been found.

APPENDIX B

FLOW CHARTS FOR FORTRAN PROGRAM EZMIN

FORTRAN program EZMIN was developed to solve the problem of finding the vector <u>r</u> which minimizes the expected cost of on hand inventory subject to a required expected number of unit systems up. This appendix contains flow charts for the main program and the following subprograms:

Subroutine PSITAB, which computes a table of values of PI(J,Y) for j=1,...,N, and Y=1,...,100, where

$$PI(J,Y) = \Psi_{J}(Y-1;r_{J} = -1)$$
.

Subroutine INITAL, which computes a set of initial values of r_j , such that $E(Z_j) \ge E(Z)_{\min}$, j = 1, ..., N.

Subroutine OPTIMR, which uses a marginal analysis technique to find an optimal r vector.

Function PRZJK, which computes $Pr\{Z_{i} = k; r_{i}\}$.

Function EBO, which computes $E(Y_{i}; r_{i})$.

Function PPT, which computes individual terms of the Poisson probability distribution.

Function PCUMT, which computes complementary cumulative terms of the Poisson probability distribution.

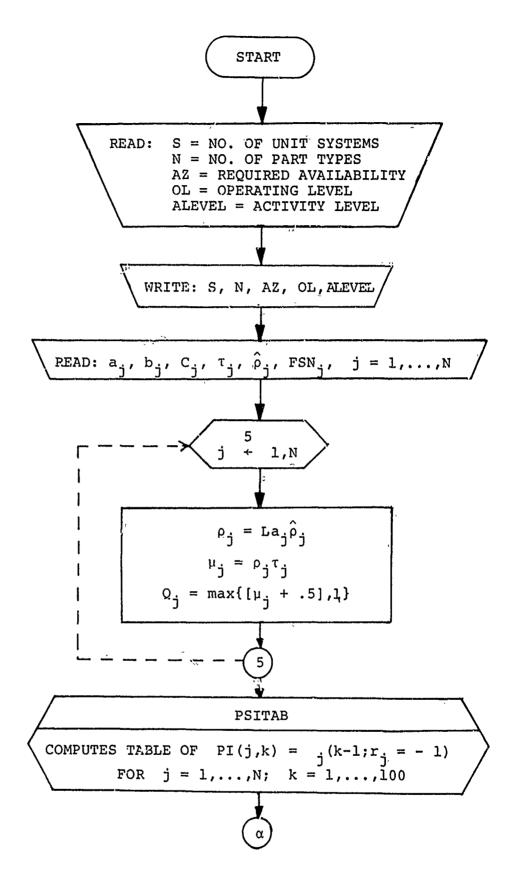
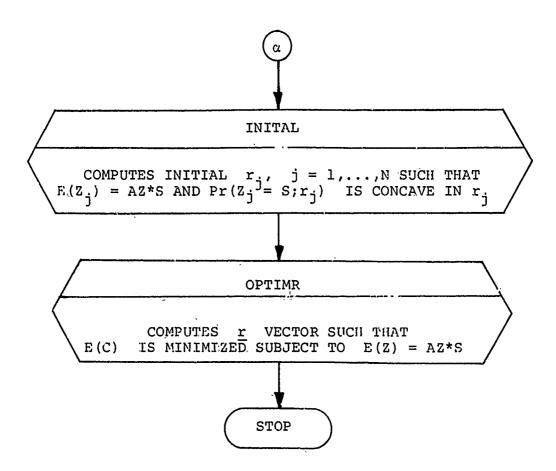


Figure Bl. Flow Chart for EZMIN Main Program



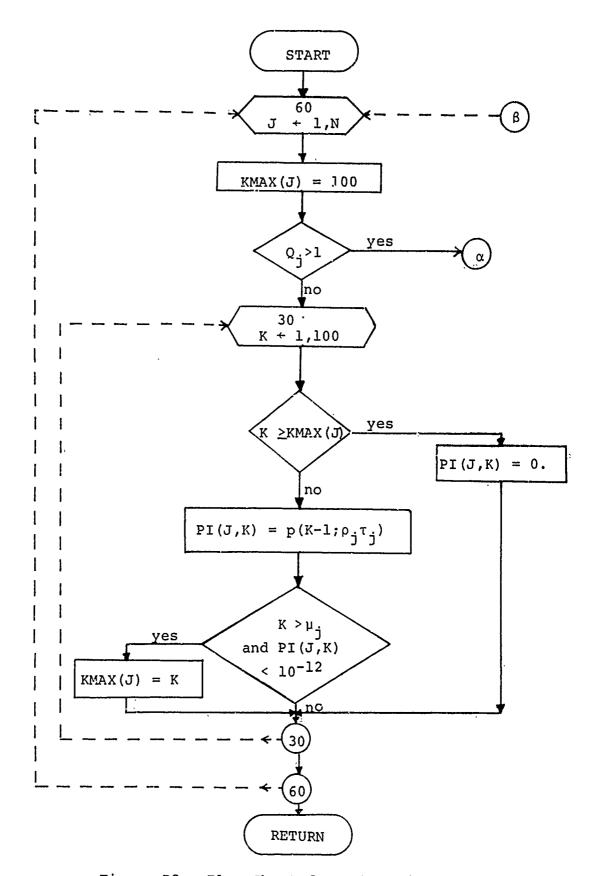


Figure B2. Flow Chart for Subroutine PSITAB

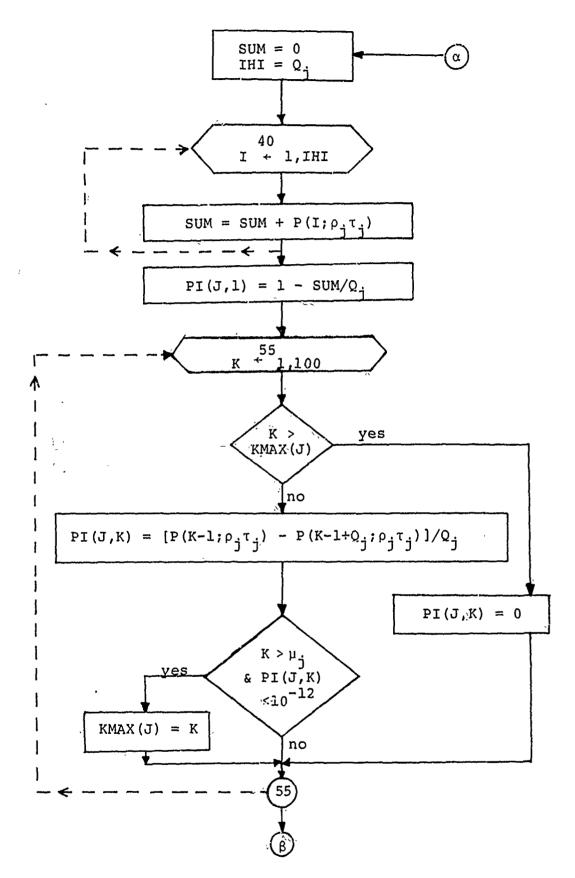


Figure B2. Continued

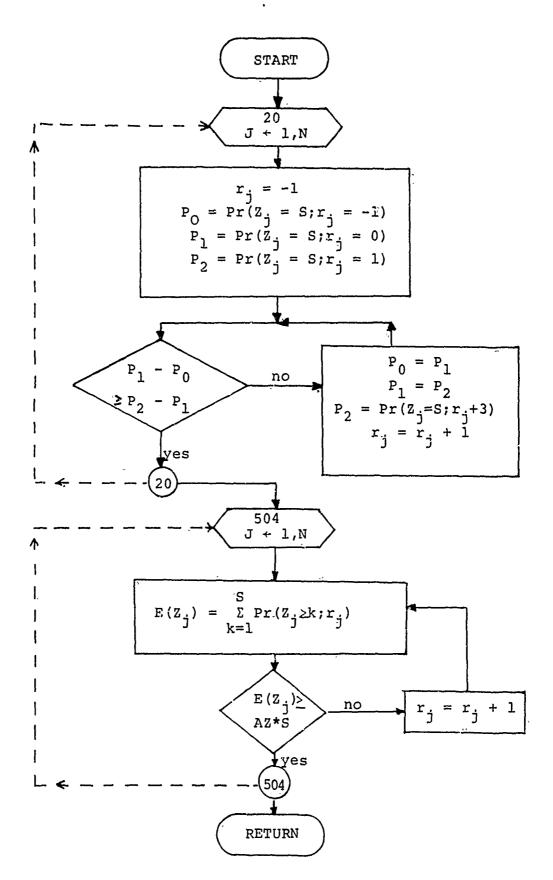


Figure B3. Flow Chart for Subroutine INITAL

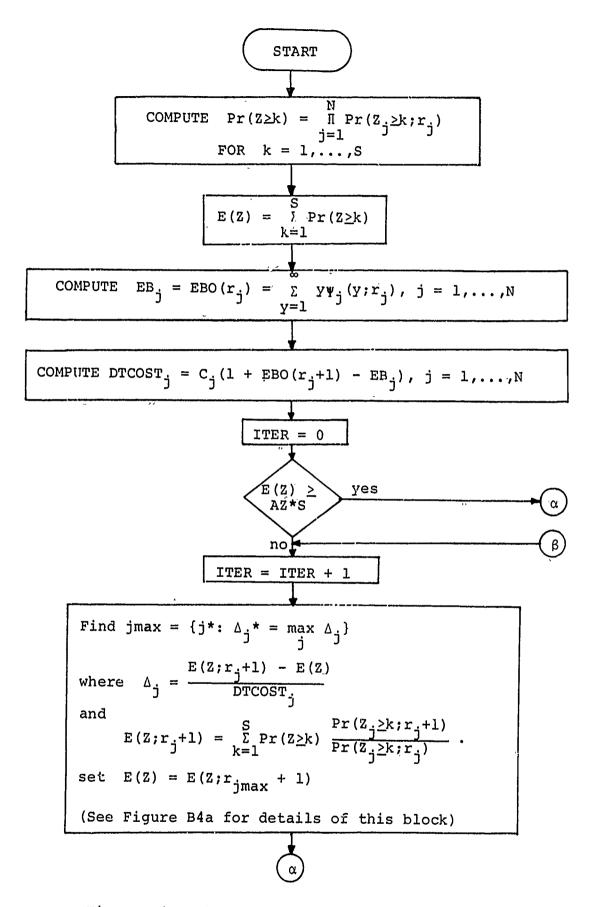


Figure B4. Flow Chart for Subroutine OPTIMR

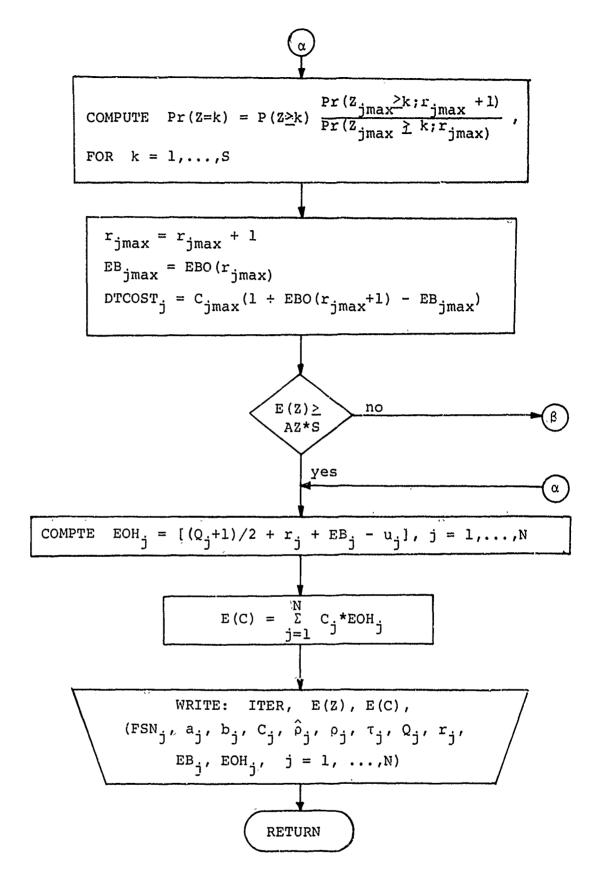


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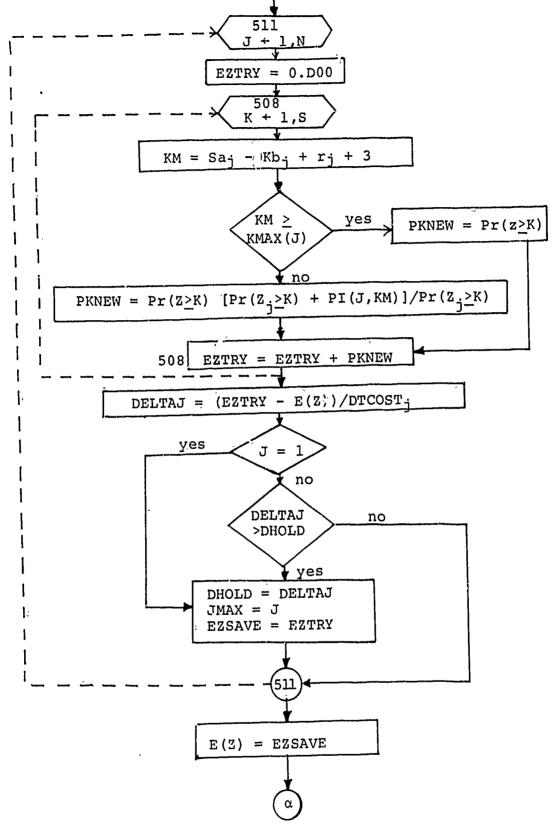


Figure B4a. Flow Chart for JMAX Segment of OPTIMR

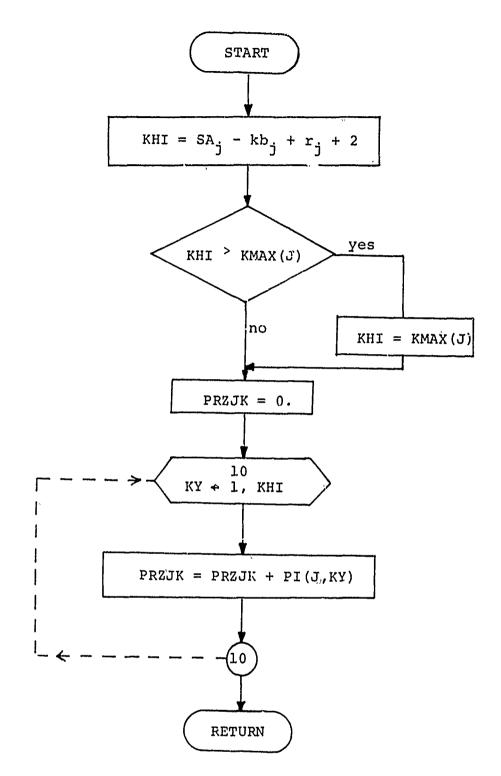


Figure B5. Flow Chart for Function PRZJK

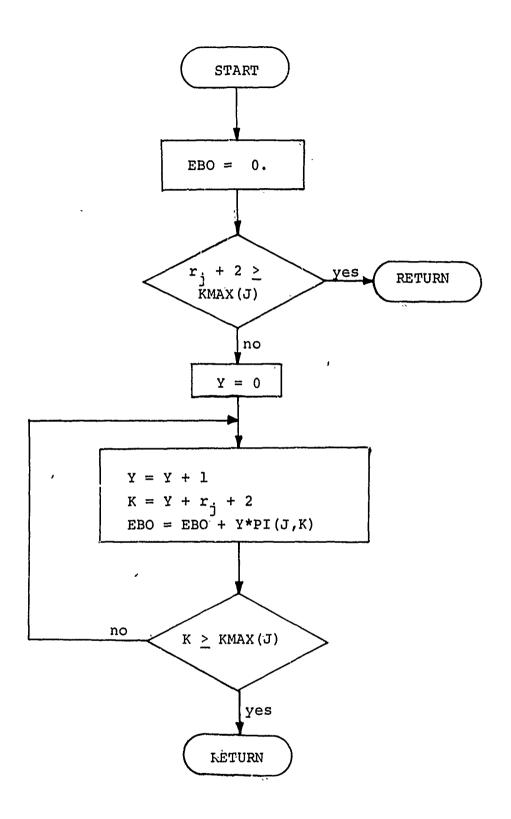


Figure B6. Flow Chart for Function EBO

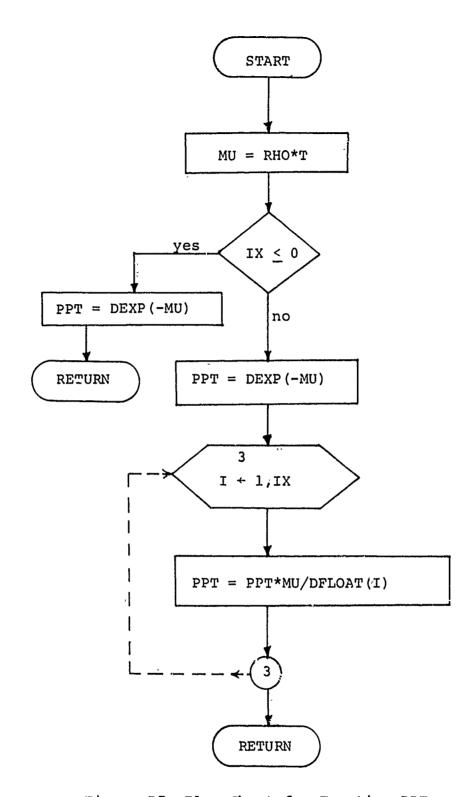


Figure B7. Flow Chart for Function PPT

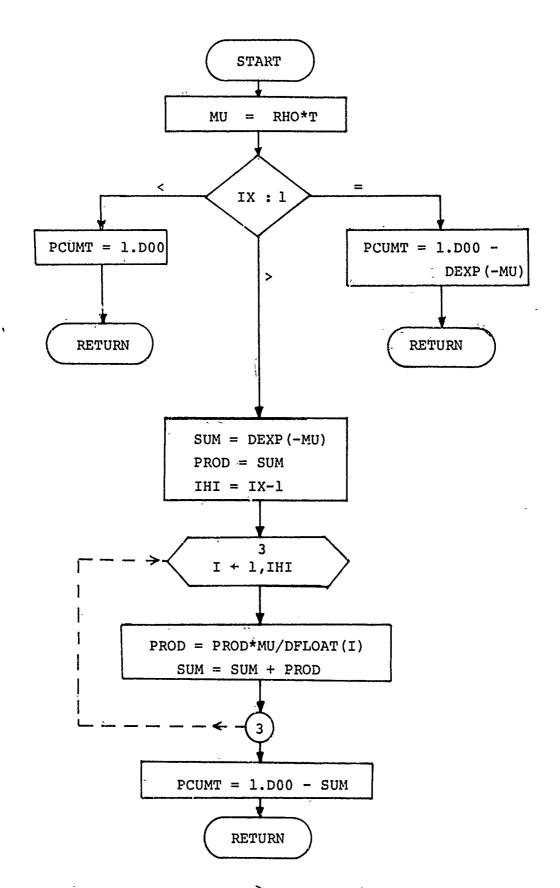


Figure B8. Flow Chart for Function PCUMT

APPENDIX C.

FLOW CHARTS FOR FORTRAN PROGRAM PKSMIN

FORTRAN program PKSMIN was developed to solve the problem of finding the <u>r</u> vector which minimizes the expected cost of on hand inventory subject to a required probability that at least k unit systems are up. This appendix contains flow charts for the main program and the following subroutines:

Subroutine INTLZ, which compute a set of initial values of r_j such that $Pr(Z_j \ge k) \ge P(Z \ge k)_{min}$, $j \ge 1, \ldots, N$.

Subroutine OPTMZ, which uses a marginal analysis technique to find an optimal \underline{r} vector.

Program PKSMIN also uses subroutine PSITAB and function subprograms PRZJK, EBO, PPT, and PCUMT, flow charts for which are in Appendix B.

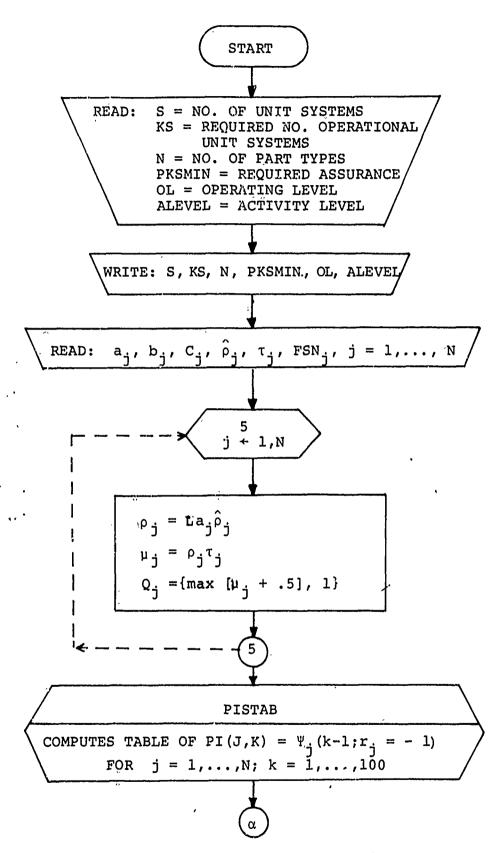


Figure Cl. Flow Chart for PKSMIN Main Program

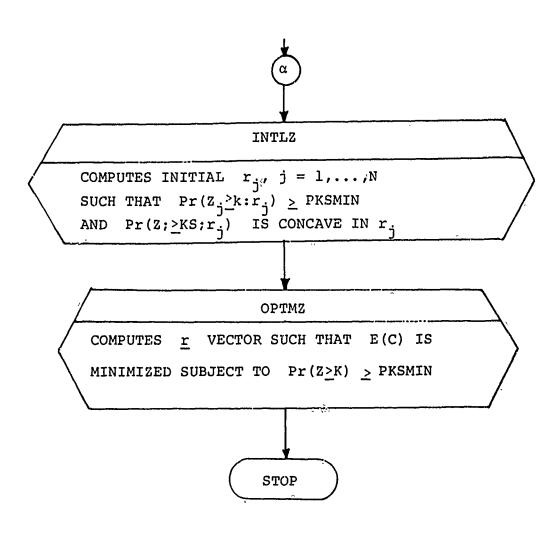


Figure Cl. Continued

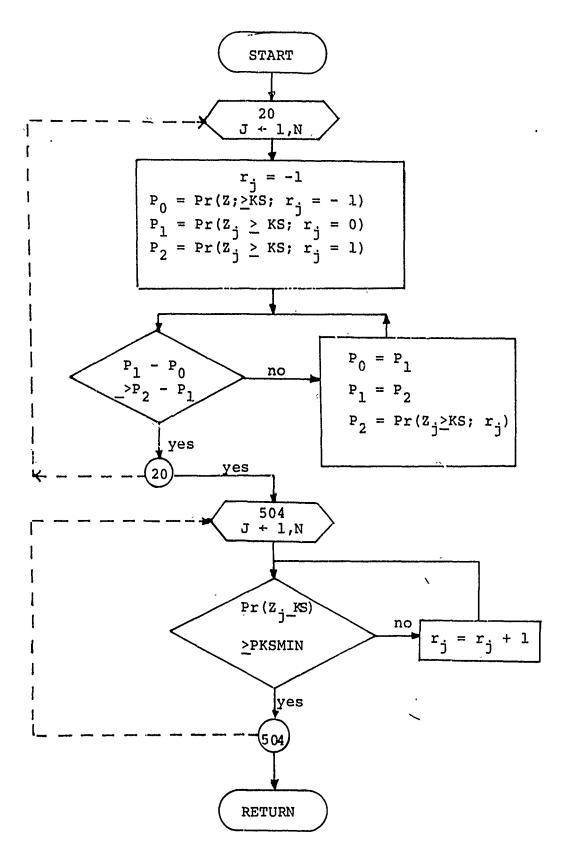


Figure C2. Flow Chart for Subroutine INTLZ

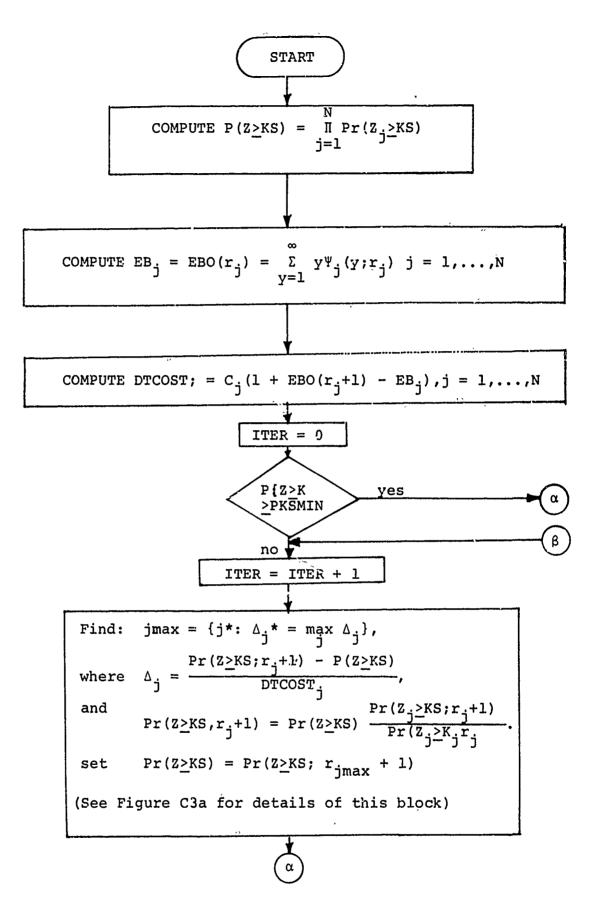
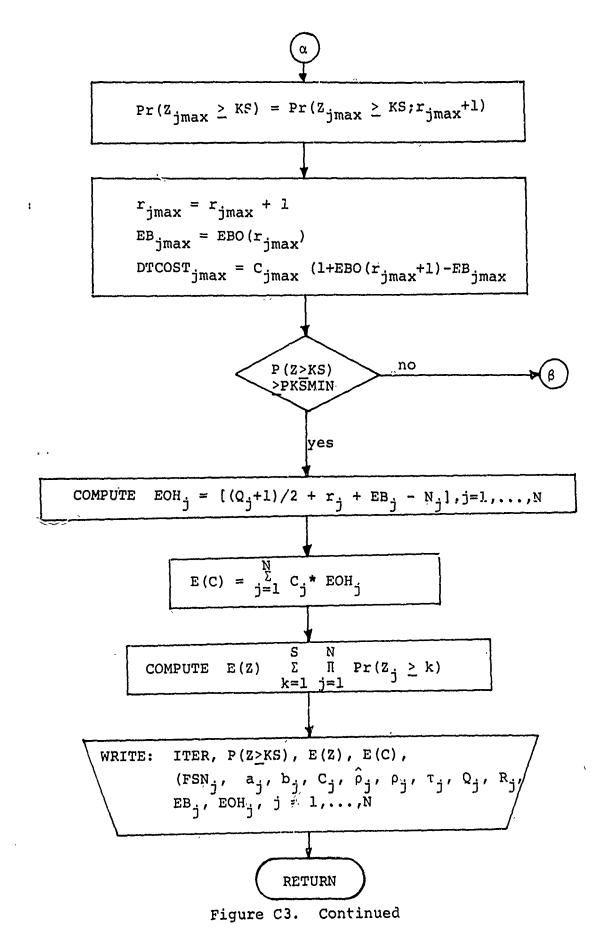


Figure C3. Flow Chart for Subroutine OPTMZ



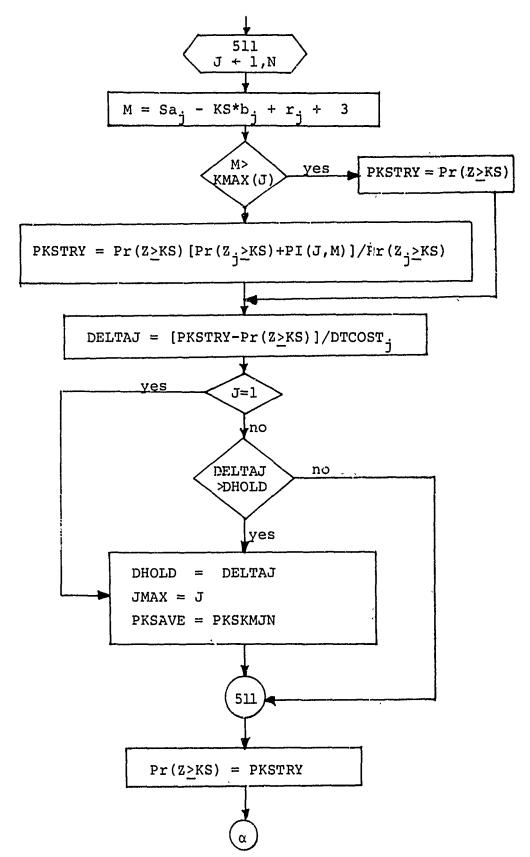


Figure C3a. Flow Chart for JAMX Segment of OPTMZ

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13. ABSTRACT				_		

A model is developed for relating continuous review inventory policies for repair parts to system availa-The system consists of S identical unit systems, each of which is a series of k-out-of-n structures. Unit system states are zero or one. An optimal cannibalization policy is assumed. Under this assumption the number of unit systems up is always the maximum possible for any given vector of backorders for the N part types in the system. The distribution of backorders under a (Q,r) policy with Poisson demands for each part type is used to derive expressions for system availability as functions of the Q,r vectors. For simplicity it is assumed that order quantities are set by an operating level in terms of days of supply. A numerical technique is presented for finding the vector of reorder points (safety levels) which minimizes the expected cost of on hand inventory subject to one of two alternative availability constraints.

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